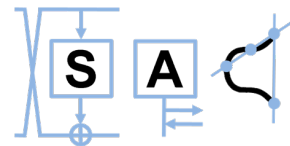


Fast and Efficient Hardware Implementation of HQC

[Sanjay Deshpande](#)^{*}, Chuanqi Xu^{*}, Mamuri Nawan⁺, Kashif Nawaz⁺, and Jakub Szefer^{*}

^{*}Computer Architecture and Security Lab, Yale University

⁺Cryptography Research Centre, Technology Innovation Institute



**Selected Areas in Cryptography
2023**



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Outline

- Motivation
- Introduction
- Goal and Existing work
- Components in HQC and their Hardware Implementation
- Comparison with Related State of the Art

Motivation

- Quantum Computers with sufficient Qubits will be able to solve practical problems.
- Two algorithms are of special interest for cryptography community:
 - Shor's Algorithm
 - Grover's Search Algorithm
- For example, to break a 2048-bit RSA, a perfect quantum computer with 4099 ideal qubits can do it in 10 seconds.
- Most recently, IBM announced 433-qubits quantum computer "Osprey".
- There is a need for Quantum-safe Cryptography!



Image Source: MIT Technology Review

Source: <https://www.quintessencelabs.com/blog/breaking-rsa-encryption-update-state-art/>

Source: <https://newsroom.ibm.com/2022-11-09-IBM-Unveils-400-Qubit-Plus-Quantum-Processor-and-Next-Generation-IBM-Quantum-System-Two>

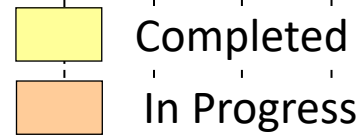
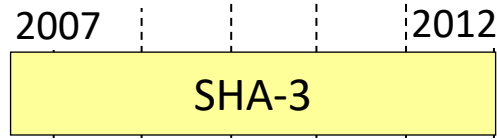


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Post Quantum Cryptography Competition

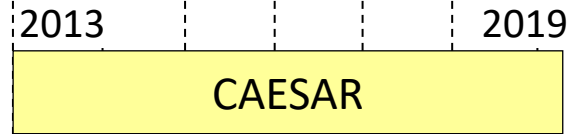
51 hash functions

→ 1 winner

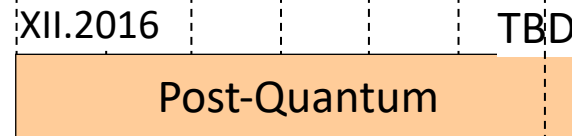


57 authenticated ciphers

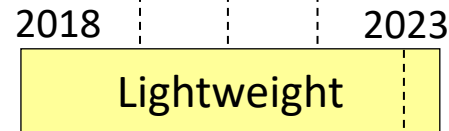
→ multiple winners



69 Public Key Post Quantum Cryptography Schemes



56 Lightweight authenticated ciphers & hash functions



07 08 09 10 11 12 13 14 15 16 17 18 19 20 21 22 23 Year

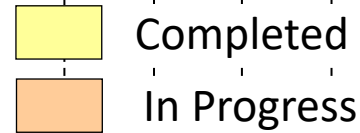
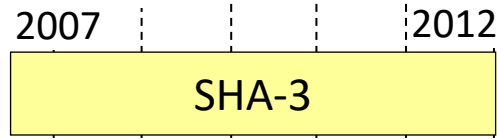


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Post Quantum Cryptography Competition

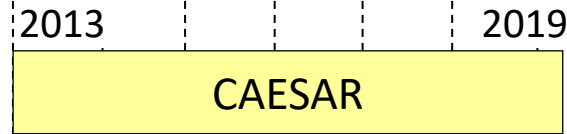
51 hash functions

→ 1 winner

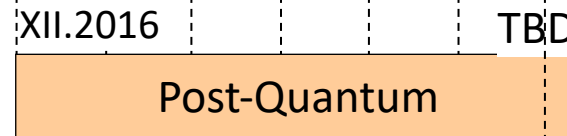


57 authenticated ciphers

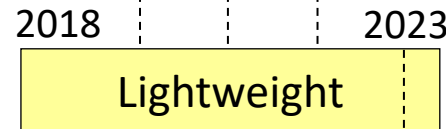
→ multiple winners



69 Public Key Post Quantum Cryptography Schemes



56 Lightweight authenticated ciphers & hash functions

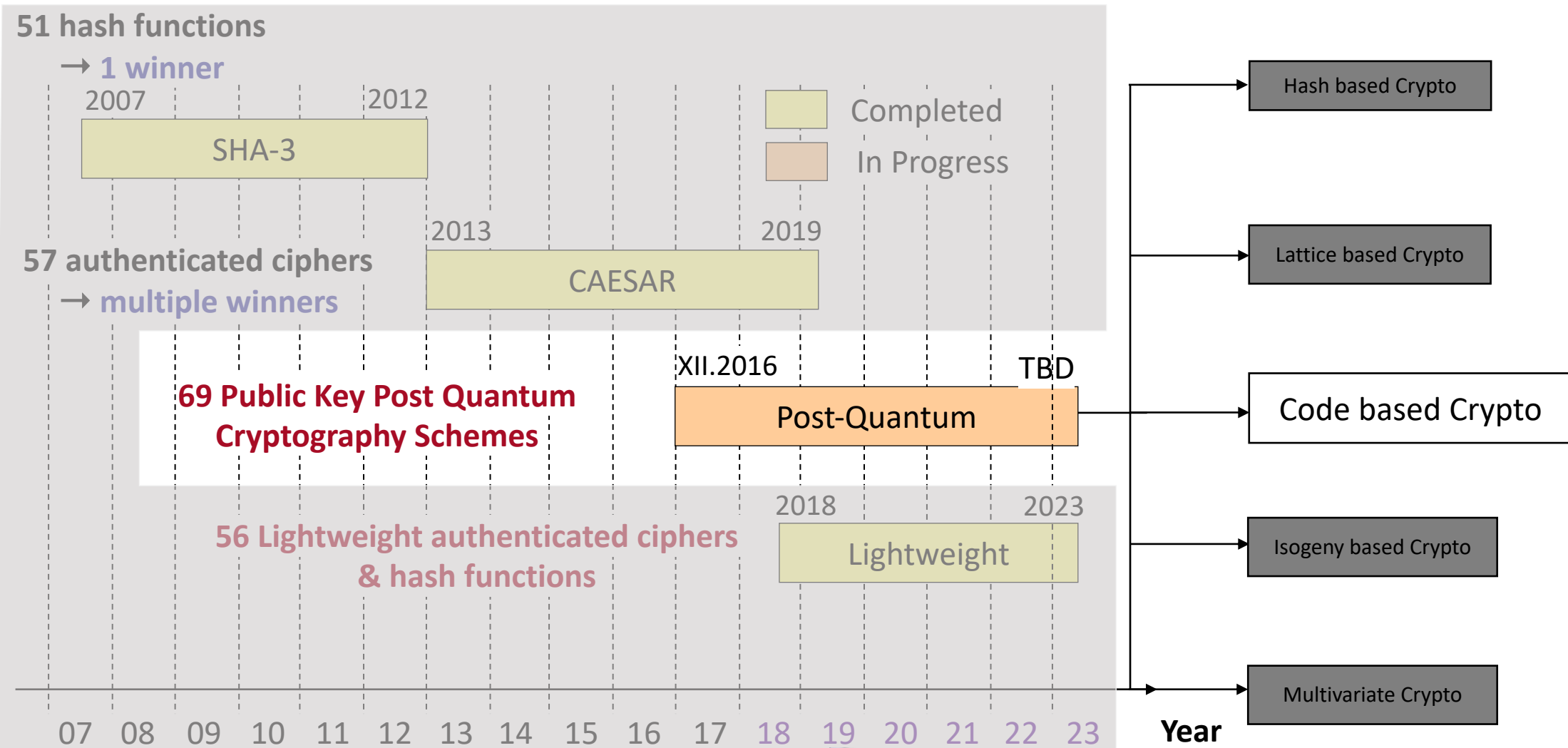


Round 4 announced
1 candidate selected for standardization
4 candidates will undergo further evaluation

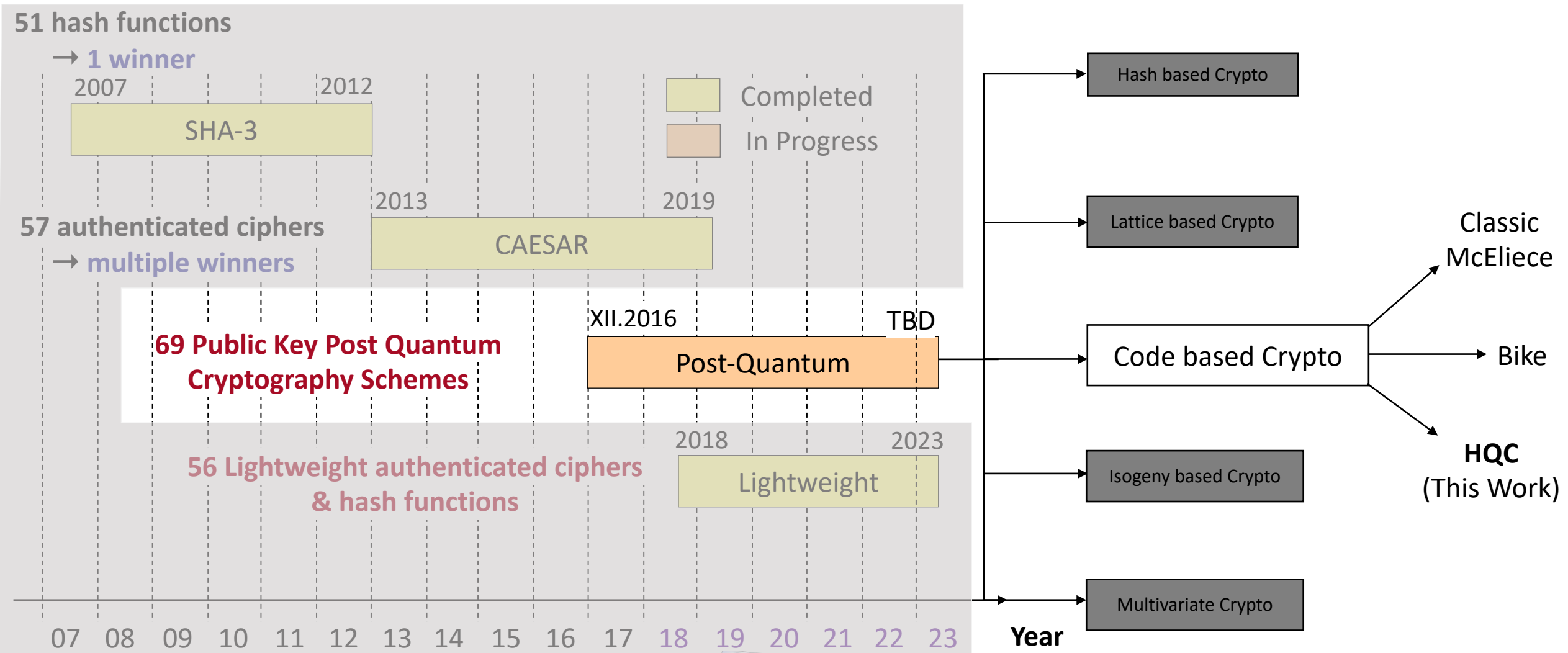
07 08 09 10 11 12 13 14 15 16 17 18 19 20 21 22 23 Year



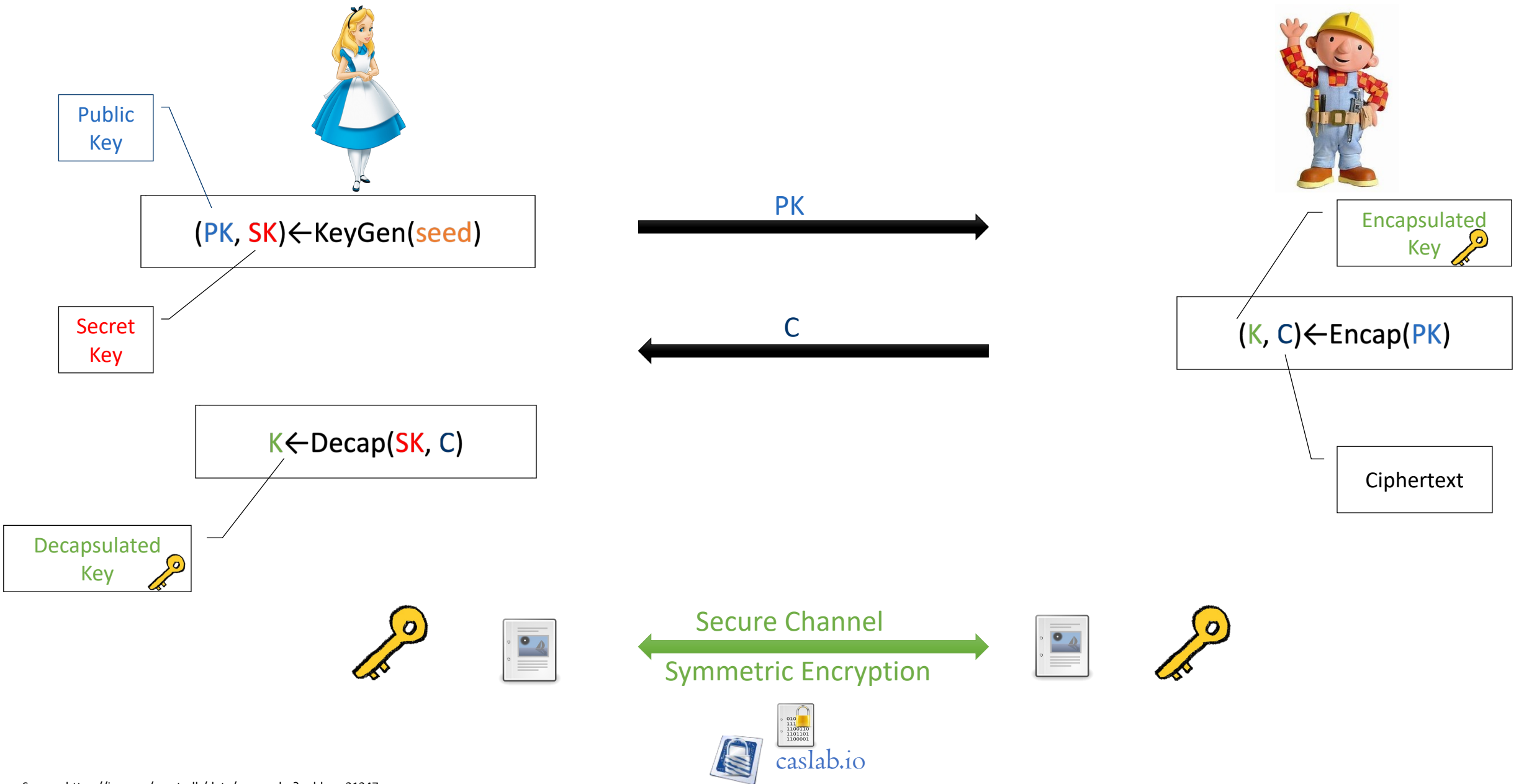
Post Quantum Cryptography Competition



Post Quantum Cryptography Competition



Key Encapsulation Mechanism



Design & Experimental Setup

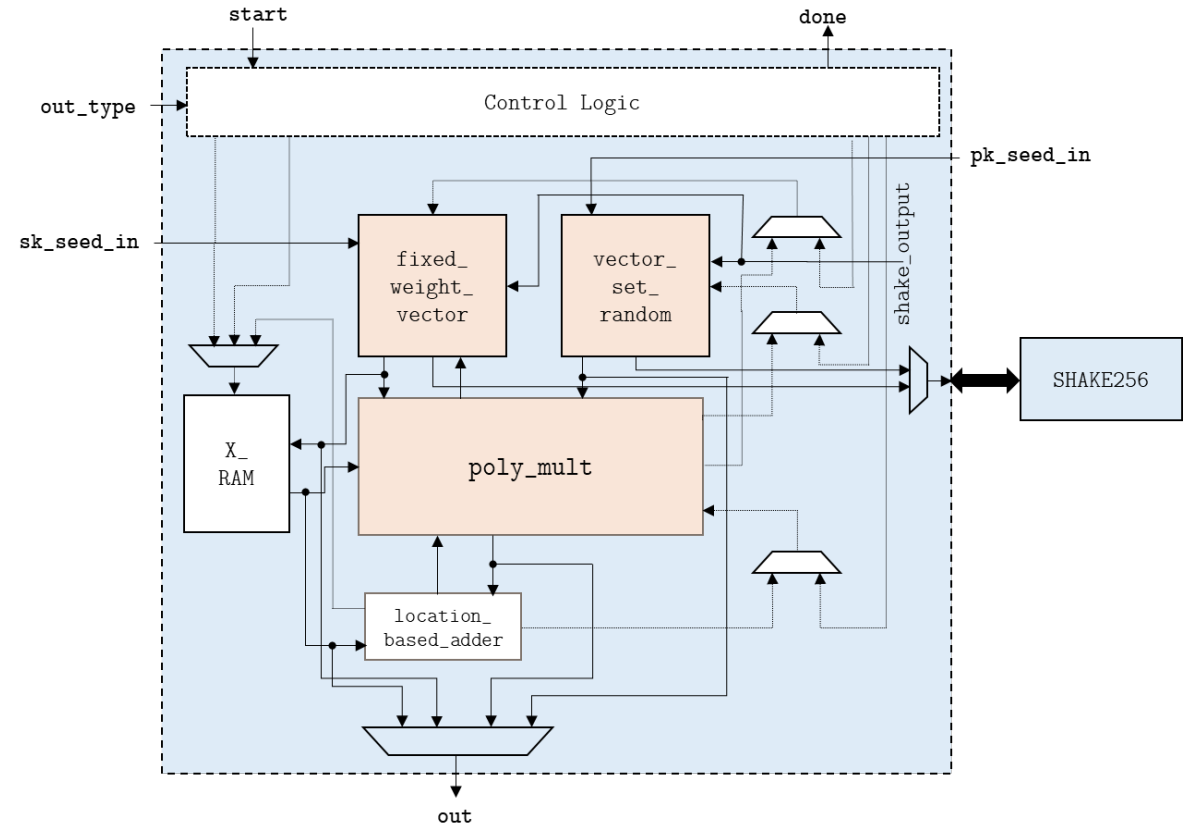
- HQC KEM parameter sets recommended as per the specification:
 - **hqc128**
 - **hqc192**
 - **hqc256**
- Target device: Xilinx - Artix 7 – xc7a200t-3 and xc7a100t-3
- Tool: Xilinx Vivado 2020.2.2 
- Verification: Using the Software Reference implementation by HQC team [AAB+20]
- Goal: Implement a constant-time hardware design which is parameterizable across
 - Security level
 - Performance parameters

Image Source: Xilinx

Primitives

1. Key Generation

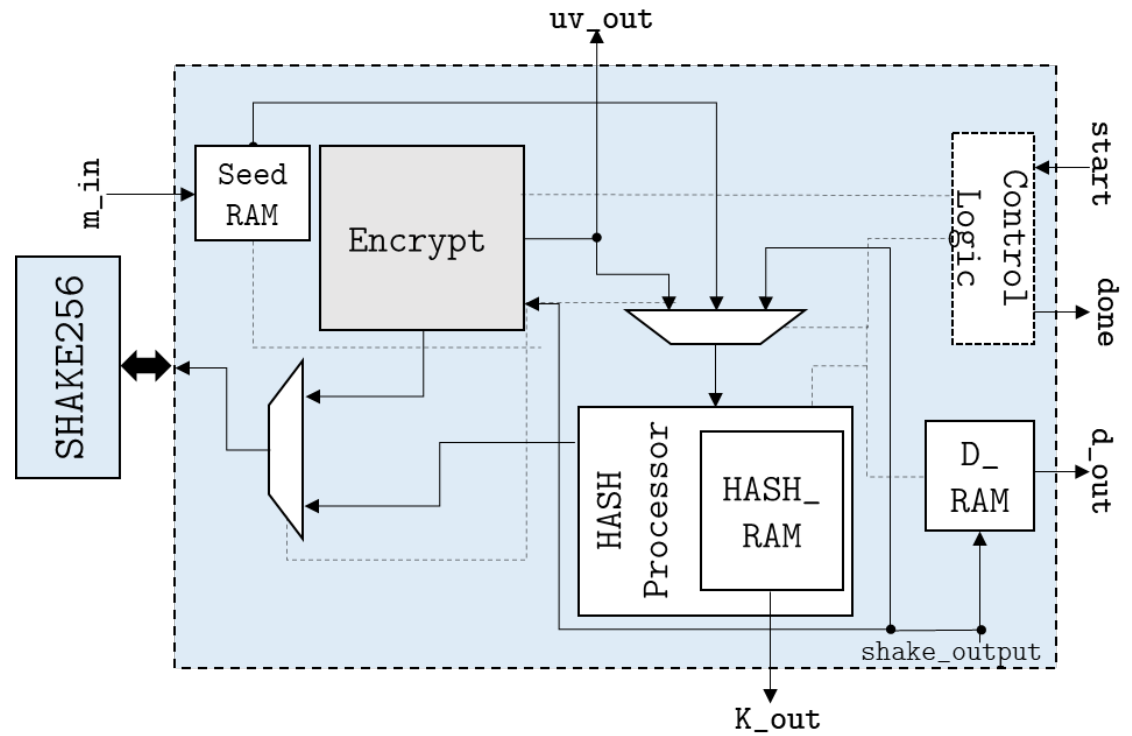
- Fixed weight error vector generation
 - SHAKE256
- Polynomial Multiplication
- Polynomial Addition



Primitives

2. Encapsulation

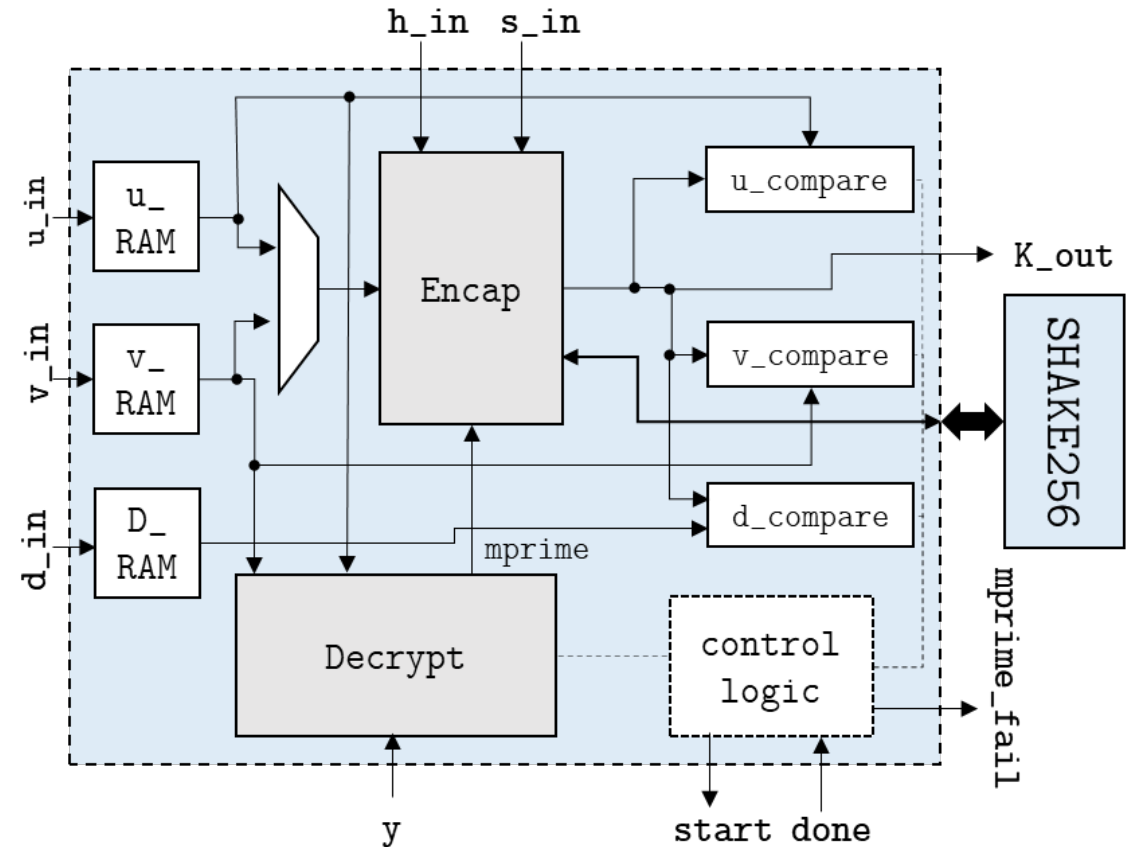
- Encrypt
 - Encode
 - Polynomial Multiplication
 - Polynomial Addition
 - Fixed weight error vector generation
 - SHAKE256
- Hash Computation
 - SHAKE256



Primitives

3. Decapsulation

- Decrypt
 - Decode
 - Polynomial Subtraction
 - Polynomial Multiplication
 - Encrypt
 - Hash Computation
 - SHAKE256
- } Encapsulation



Primitives – Common Modules

1. Key Generation

- Fixed weight error vector generation
 - SHAKE256
- Polynomial Multiplication
- Polynomial Addition

3. Decapsulation

- Decrypt
 - Decode
 - Polynomial Subtraction
 - Polynomial Multiplication
- Encrypt
- Hash Computation
 - SHAKE256

2. Encapsulation

- Encrypt
 - Encode
 - Polynomial Multiplication
 - Polynomial Addition
 - Fixed weight error vector generation
 - SHAKE256
- Hash Computation
 - SHAKE256

Primitives - Existing Work

1. Key Generation

- Fixed weight error vector generation
 - SHAKE256 [CCD+22]
- Polynomial Multiplication
- Polynomial Addition

3. Decapsulation

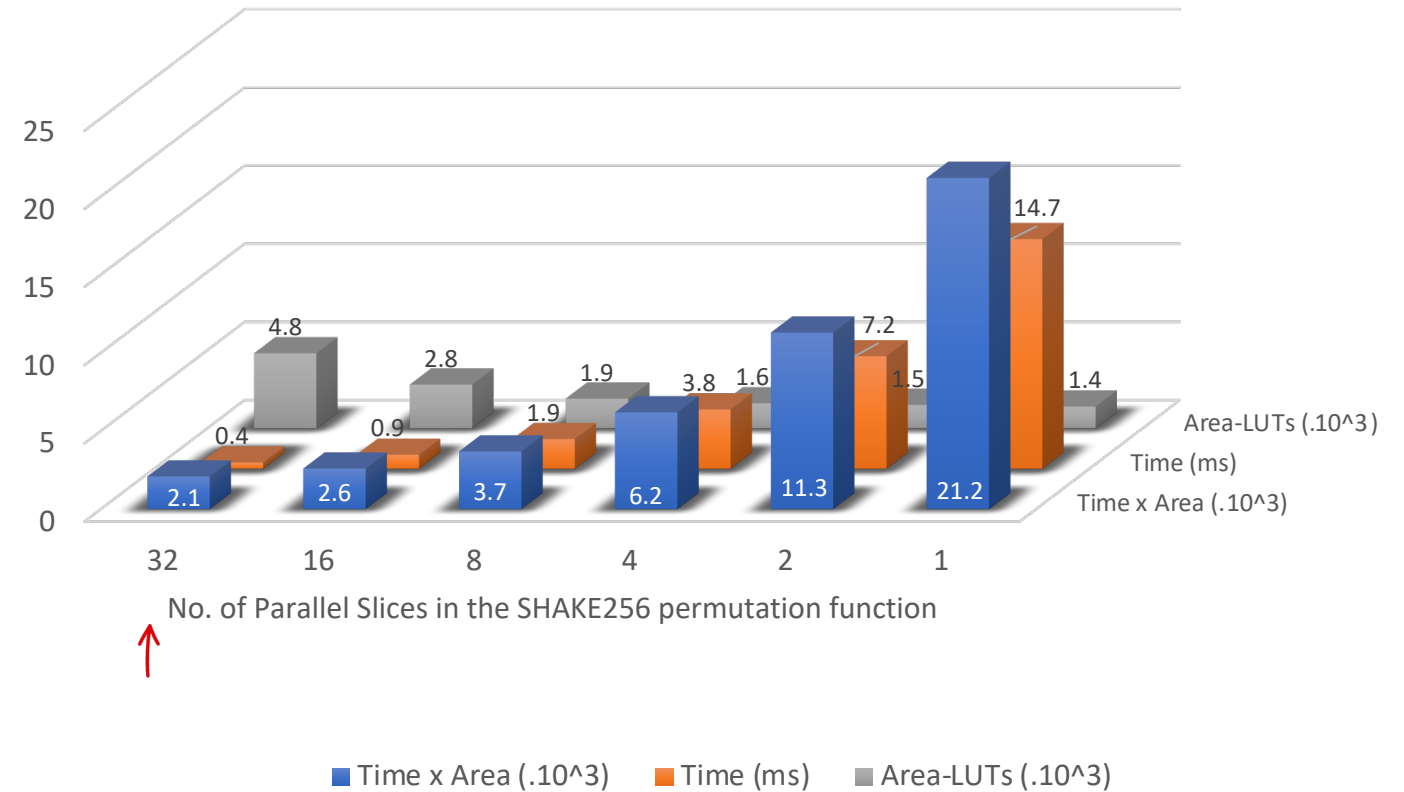
- Decrypt
 - Decode
 - Polynomial Subtraction
 - Polynomial Multiplication
- Encapsulation

2. Encapsulation

- Encrypt
 - Encode
 - Polynomial Multiplication
 - Polynomial Addition
 - Fixed weight error vector generation
 - SHAKE256 [CCD+22]
- Hash Computation
 - SHAKE256 [CCD+22]

SHAKE256

- Improved existing design [CCD+22].
 - Added state preserving capability.
 - Added an extra mode in the performance parameter (Parallel Slices).



Hardware Implementation of HQC KEM

1. Key Generation

- Fixed weight error vector generation
 - SHAKE256
- Polynomial Multiplication
- Polynomial Addition

3. Decapsulation

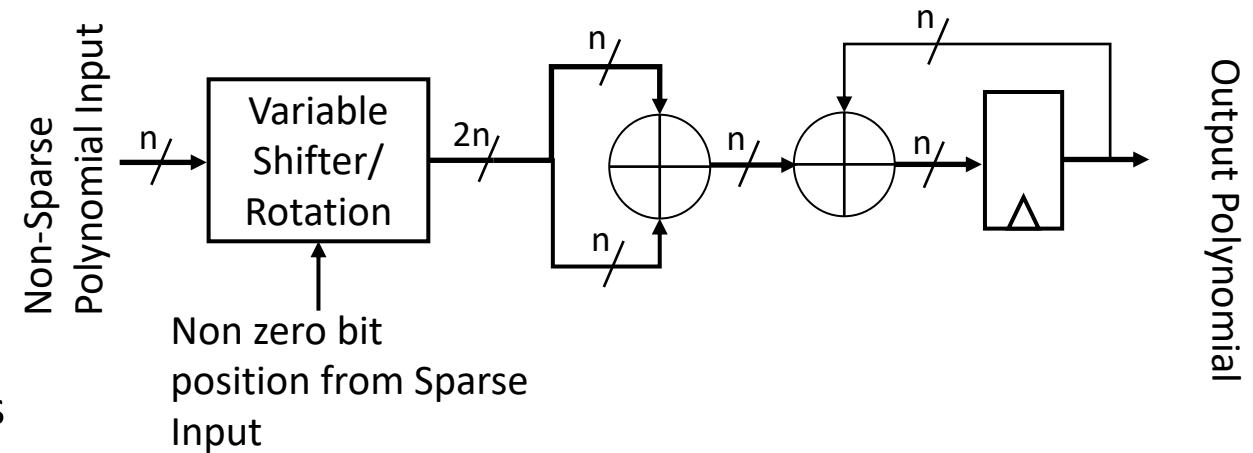
- Decrypt
 - Decode
 - Polynomial Subtraction
 - Polynomial Multiplication
- Encapsulation

2. Encapsulation

- Encrypt
 - Encode
 - Polynomial Multiplication
 - Polynomial Addition
 - Fixed weight error vector generation
 - SHAKE256
- Hash Computation
 - SHAKE256

Polynomial Multiplication - Sparse Multiplication with Interleaved Reduction

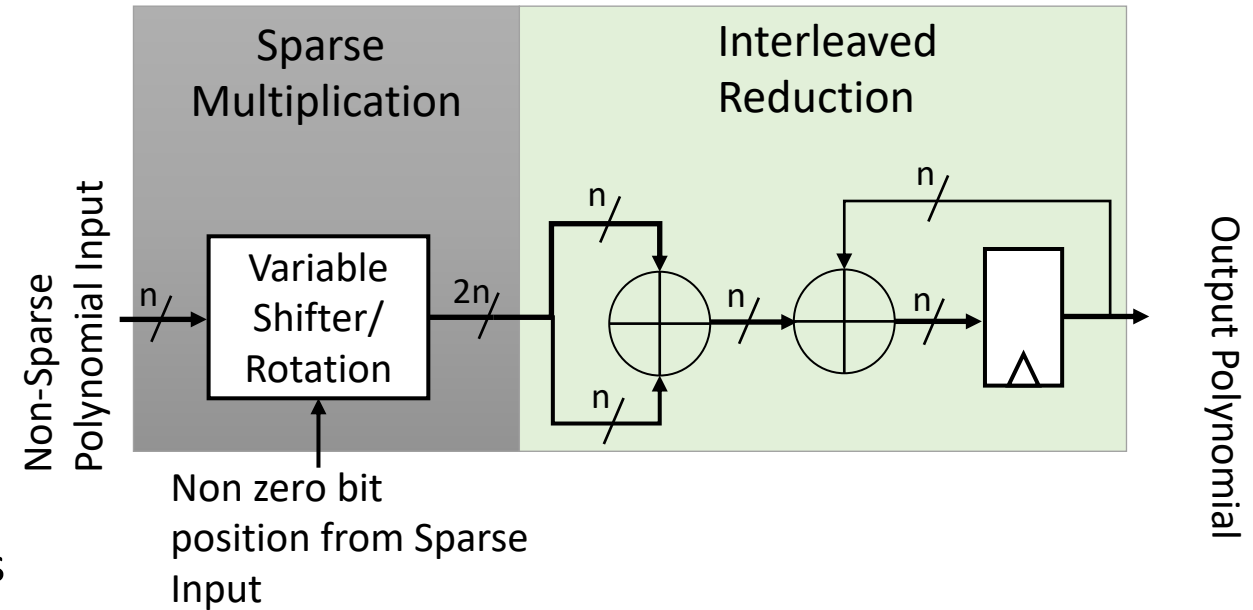
- One of the inputs to polynomial multiplier is a sparse fixed weight vector.
- Indices of non-zero elements are used to shift non-sparse polynomial to imitate the multiplication.
- The multiplication and the modular reduction is interleaved.



*Conceptual Block Diagram

Polynomial Multiplication - Sparse Multiplication with Interleaved Reduction

- One of the inputs to polynomial multiplier is a sparse fixed weight vector.
- Indices of non-zero elements are used to shift non-sparse polynomial to imitate the multiplication.
- The multiplication and the modular reduction is interleaved.



*Conceptual Block Diagram

Polynomial Multiplication - Results

BW (bits)	Resources						F (MHz)	Cycles (cyc.)	Time (us)	T x A
	Logic			Memory						
	(SLICES)	(LUT)	(DSP)	(FF)	(BR)					
Our poly_mult module, Polynomial Length ⁺ = 12,323, $W_{SPARSE}^* = 71$										
32	134	396	0	181	1	270	27,621	0.10	14	
64	202	599	0	205	2	277	13,918	0.05	10	
128	486	1,438	0	456	4	238	7,102	0.03	14	
General Sparse Multiplier, Polynomial Length ⁺ = 12,323, $W_{SPARSE}^* = 71$ [19]										
32	132	319	0	127	2	234	27,691	0.12	16	
64	197	549	0	190	4	222	13,988	0.06	12	
128	378	1,136	0	381	8	185	7,172	0.04	15	
Sparse Multiplier, Polynomial Length ⁺ = 10,163, $W_{SPARSE}^* = 71$ [15]										
32	100		—	—	2	240	158,614	0.66	66	
64	157		—	—	3	220	90,880	0.41	64	
128	292		—	—	5	210	51,688	0.24	70	

[†] = Slices (no info on LUTs), ⁺ Length of the non-sparse arbitrary polynomial, ^{*} = Weight of the sparse polynomial input

Target Device: Artix 7 – xc7a200t

Fixed Weight Vector Generation

- Generates a uniform random n -bit fixed-weight (w) vector.
- Algorithm given in [AAB+20].

Input: N , w , seed

Output: w distinct elements in range 0 to $N - 1$

1: $pos = [] * w$

2: $i, j, k, l = 0$

3: $prng_init(seed)$

4: while ($i < w$) do

5: $rand_bits \leftarrow prng_draw(outsize = 24)$

6: if ($rand_bits < THRESHOLD_VALUE$)

7: $pos[i] = (rand_bits \% N)$

8: $i++$

9: end

10: end while

Threshold Check

```
11: while (j < w) do
12:     while (k < w) do
13:         if (pos[j] == pos[k]) then
14:             while (l < 1) do
15:                 rand_bits ← prng_draw(outsize=24)
16:                 if (rand_bits < THRESHOLD) then
17:                     pos[k] = (rand_bits % N)
18:                     l++
19:                 end if
20:             end while
21:         end
22:         else then
23:             k++
24:         end
25:     end while
26:     j++
27: end while
28: return pos
```

Duplicate Detection

//Value for THRESHOLD is a constant based on the parameter set



Fixed Weight Vector – Constant Weight Word (CWW) Method

- Constant time method proposed by Sendrier [SEN21].
- Recommended by HQC Team as a 4th round update.
- Barrett Reduction used in Index Generation.
- However, this method has small bias.

Input: N, w, seed

Output: w distinct elements in range 0 to $N - 1$

```
1: rand bits  $\leftarrow$  prng(input = seed, output size =  $32 \times w$ )
2: for  $i \leftarrow w - 1$  to  $0$  do
3:     pos[i] =  $i + (\text{rand bits}[32 + 32 * i - 1 : 32 * i]) \% (N - i)$ 
4: end for
5: for  $j \leftarrow w - 1$  to  $0$  do
6:     duplicate found  $\leftarrow 0$ 
7:     for  $k \leftarrow j + 1$  to  $w - 1$  do
8:         if pos[j] == pos[k] then
9:             duplicate found  $\leftarrow 1$ 
10:        end if
11:    end for
12:    if duplicate found == 1 then
13:        pos[j] =  $j$ 
14:    end
15: end for
16: return pos
```

RNG

Index Generation

Duplicate Detection and Index Replacement



Fixed Weight Vector Generation – Fast and Non-Biased(FNB) Method

- Improvement on the original algorithm proposed given in [AAB+20]

Input: N, w, seed, ACC_REJ

Output: w distinct elements in range 0 to N – 1

1: pos = []*(w + ACC_REJ)

2. i,j,k,m = 0

3. prng_init(seed)

4: while (i < (w+ ACC_REJ)) do

5: rand_bits ← prng_draw(outsize = 24)

6: if (rand_bits < THRESHOLD_VALUE) then

7: pos [i] = (rand bits % N)

8: i++

9: end

10: end while

//Value for THRESHOLD is a constant based on the parameter set

// ACC_REJ = ACCEPTABLE_REJECTIONS

// e.g., ACC_REJ = 75

```

11: while (j < w + ACC_REJ) do
12:     while (k < (w+ ACC_REJ)) do
13:         if (j < w) then
14:             if pos[j] == pos[k] then
15:                 if (m < ACC_REJ) then
16:                     pos[k] = pos[m+w]
17:                     m++
18:                 end
19:             else
20:                 USE PRNG TO DRAW MORE RAND BITS AND
21:                 REDO THRESHOLD CHECK and ASSIGN TO pos[k]
22:             end
23:         end
24:         DUMMY OPERATIONS
25:     end
26:     else
27:         DUMMY OPERATIONS
28:     end
29:     k++
30: end while
31: j++
32: end while
33: return pos
    
```

Small prob non-constant
time behavior

Duplicate Detection

Fixed Weight Vector Generation - Evaluation

Design	Resources								T x A	Failure ⁺ Prob.
	Weight (w_r)	Logic (LUT)	(DSP)	Memory (FF)	(BR)	F (MHz)	Cycles (cyc.)	Time (us)		
Fast and Non-Biased Design (ACCEPTABLE_REJECTIONS = w_r)										
hqc128	75	316	0	124	2.0	223	1,479	6.63	2.10	2.8×2^{-199}
hqc192	114	295	0	125	2.0	236	2,226	9.43	2.78	1.1×2^{-280}
hqc256	149	314	0	192	2.5	242	3,248	13.42	4.21	4.9×2^{-355}
Constant Weight Word (CWW)										
hqc128	75	201	4	229	1.0	201	3,062	15.23	3.06	0
hqc192	114	211	5	245	1.0	200	6,817	34.09	7.19	0
hqc256	149	216	5	248	1.0	204	11,487	56.31	12.6	0

⁺ = Probability of our design failing to behave constant-time.

Key Generation - Algorithm and Hardware Design

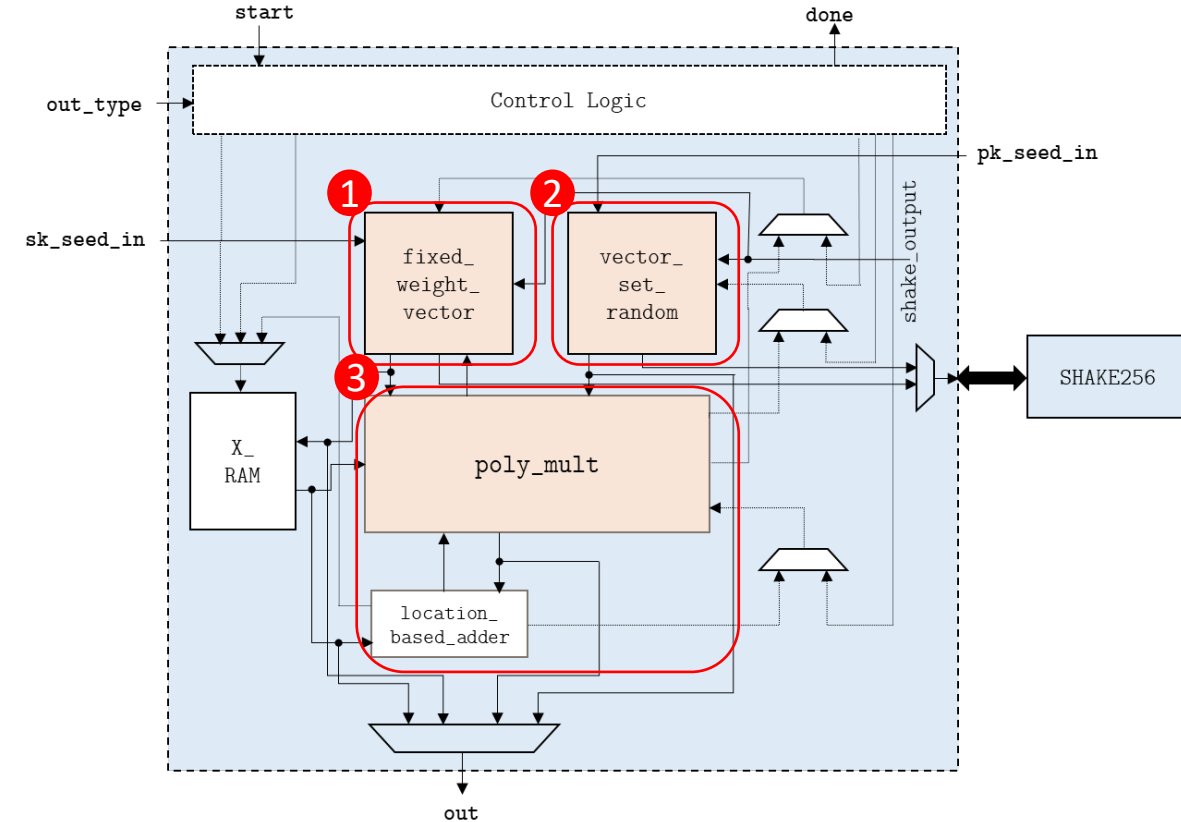
Inputs: pk_seed, sk_seed

Outputs: public key(h, s), secret key (x, y)

1. $(x, y) = \text{FixedWeight}(\text{sk_seed})$

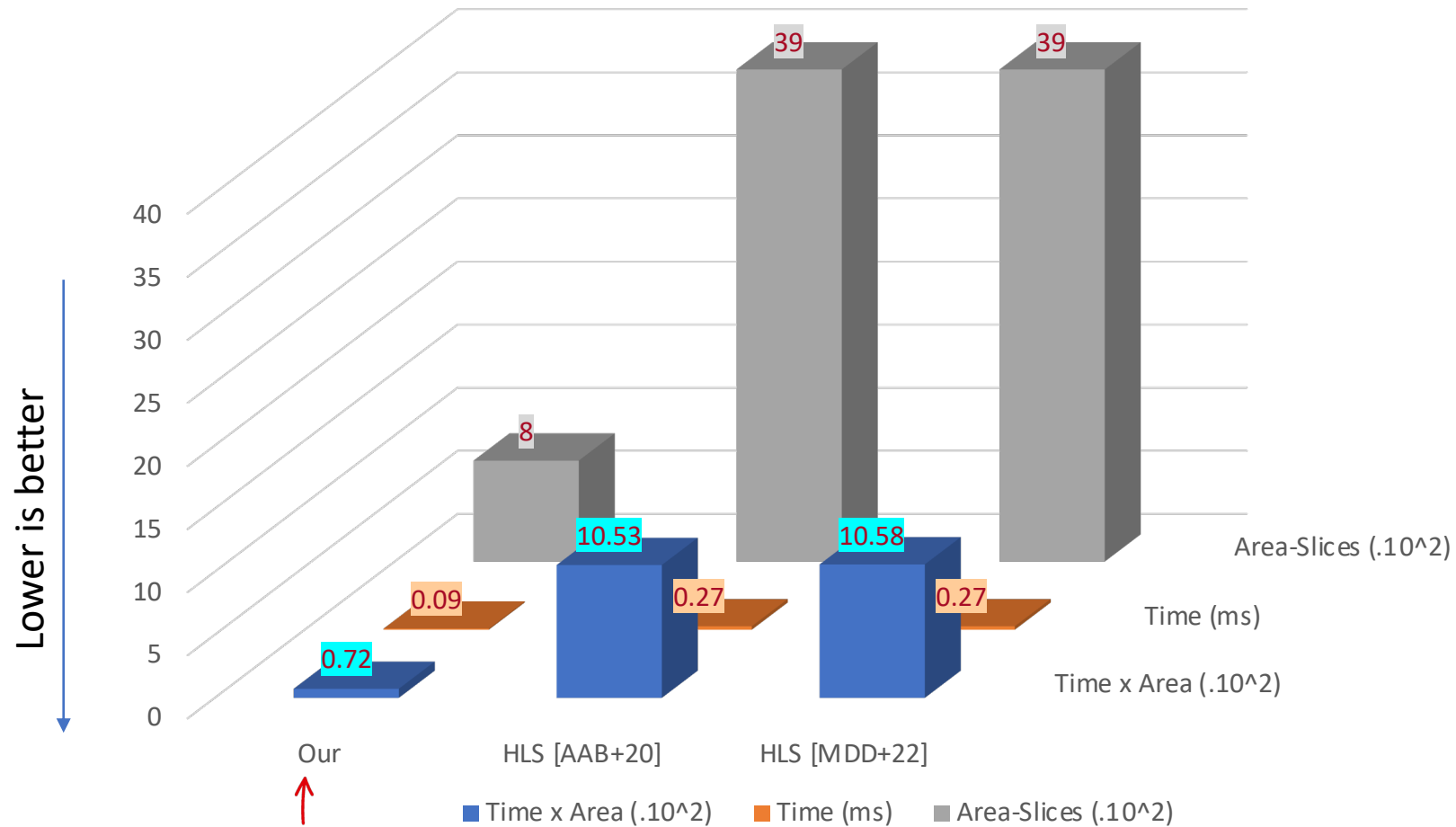
2. $h = \text{VectorRandom}(\text{pk_seed})$

3. $s = x + h.y$



Key Generation – Hardware Design Performance Comparison

HQC128



Hardware Implementation of HQC KEM

1. Key Generation

- Fixed weight error vector generation
 - SHAKE256
- Polynomial Multiplication
- Polynomial Addition

3. Decapsulation

- Decrypt
 - Decode
 - Polynomial Subtraction
 - Polynomial Multiplication
- Encapsulation

2. Encapsulation

- Encrypt
 - Encode
 - Polynomial Multiplication
 - Polynomial Addition
 - Fixed weight error vector generation
 - SHAKE256
- Hash Computation
 - SHAKE256

Encrypt - Algorithm and Hardware Design

Inputs: theta, Public Key (h_in, s_in),
message (m_in)

Outputs: Ciphertext (u_out, v_out)

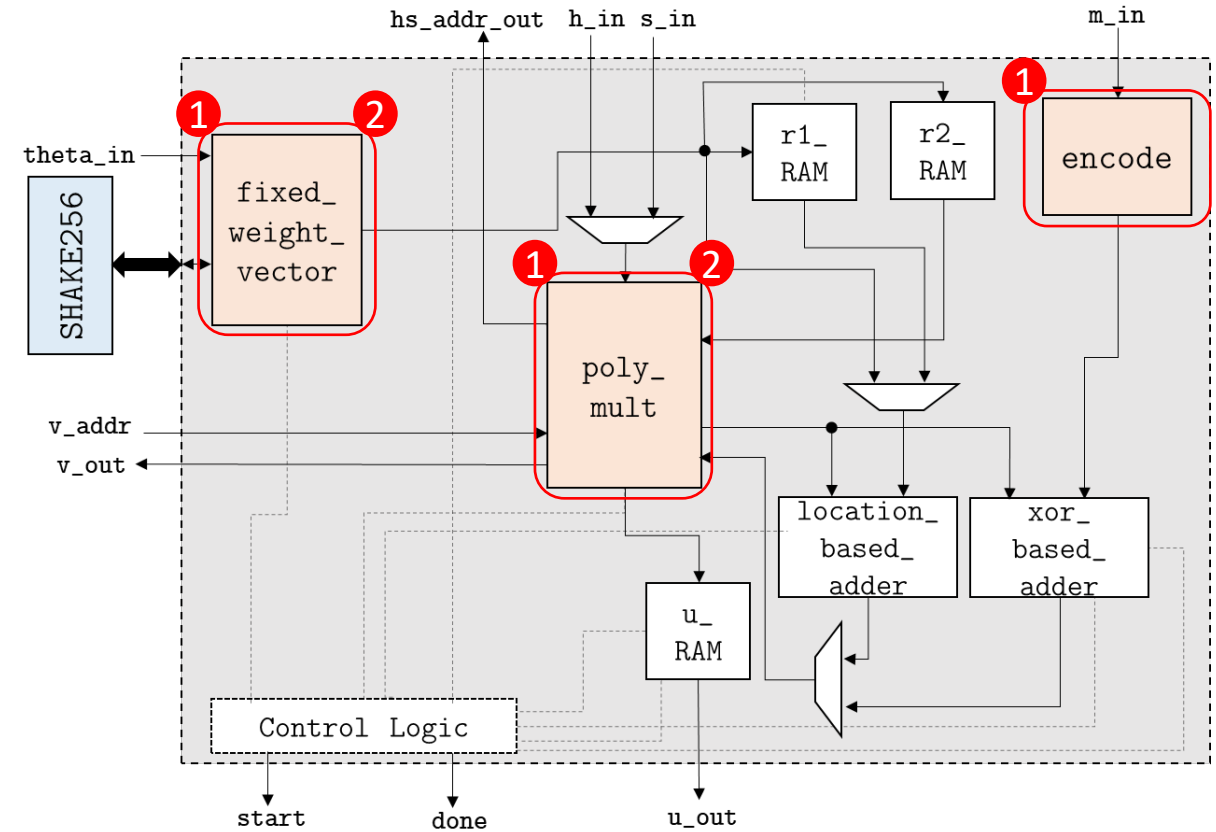
1. $u_out = r_1 + h.r_2$

2. $v_out = m_in.G + S.r_2 + e$

Where m_inG = Encoded message

r_1, r_2, e are fixed weight vectors

- Encode consists of Reed-Solomon and Reed-Muller Encoding



Encrypt - Algorithm and Hardware Design

Inputs: θ , Public Key (h_{in} , s_{in}),
message (m_{in})

Outputs: Ciphertext (u_{out} , v_{out})

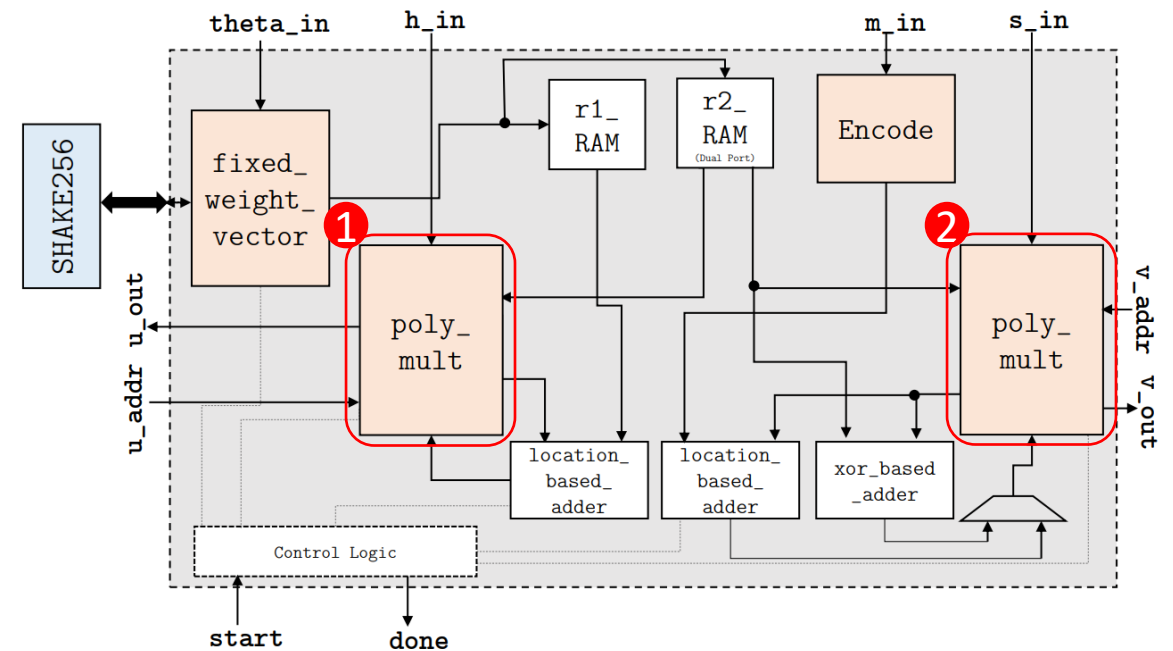
$$1. \quad u_{out} = r_1 + h.r_2$$

$$2. \quad v_{out} = m_{in}.G + S.r_2 + e$$

Where $m_{in}G$ = Encoded message

r_1 , r_2 , e are fixed weight vectors

- Two polynomial multiplications can be run in parallel.

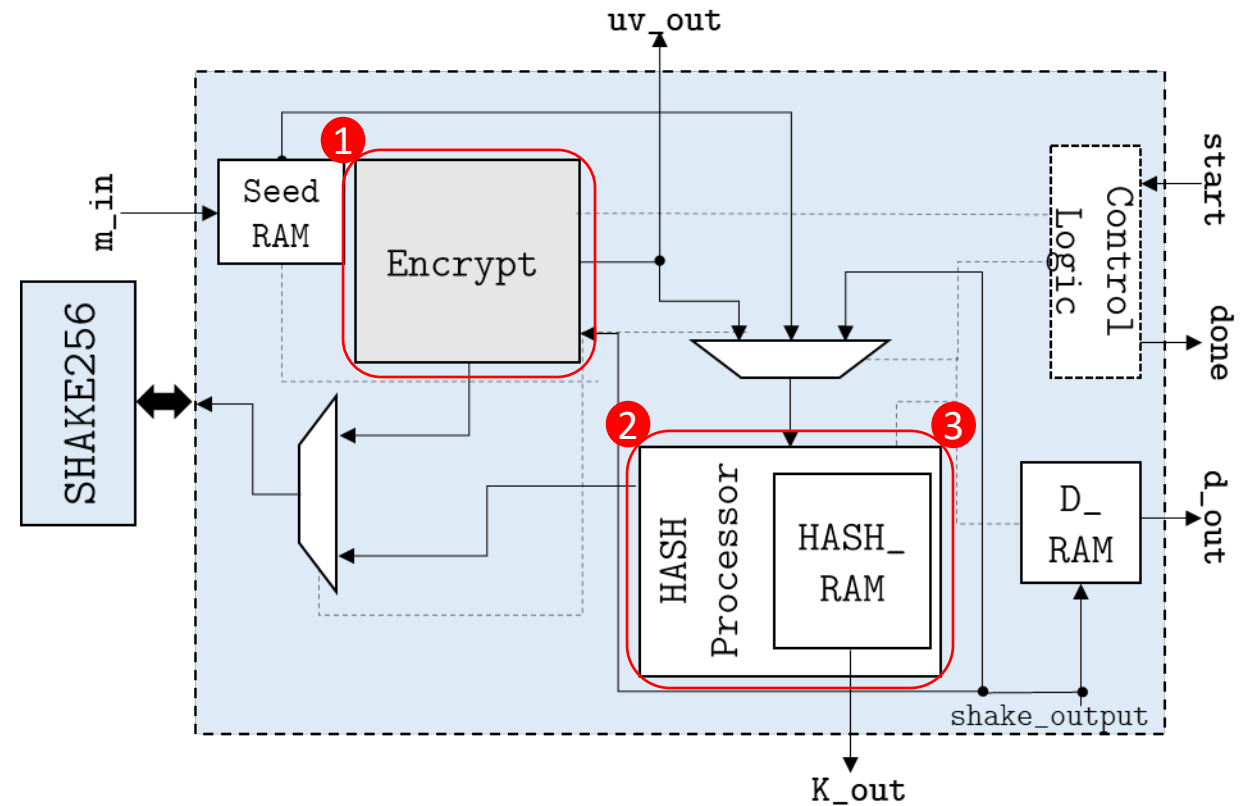


Encapsulation - Algorithm and Hardware Design

Inputs: Public Key (h_{in} , s_{in}), message (m_{in})

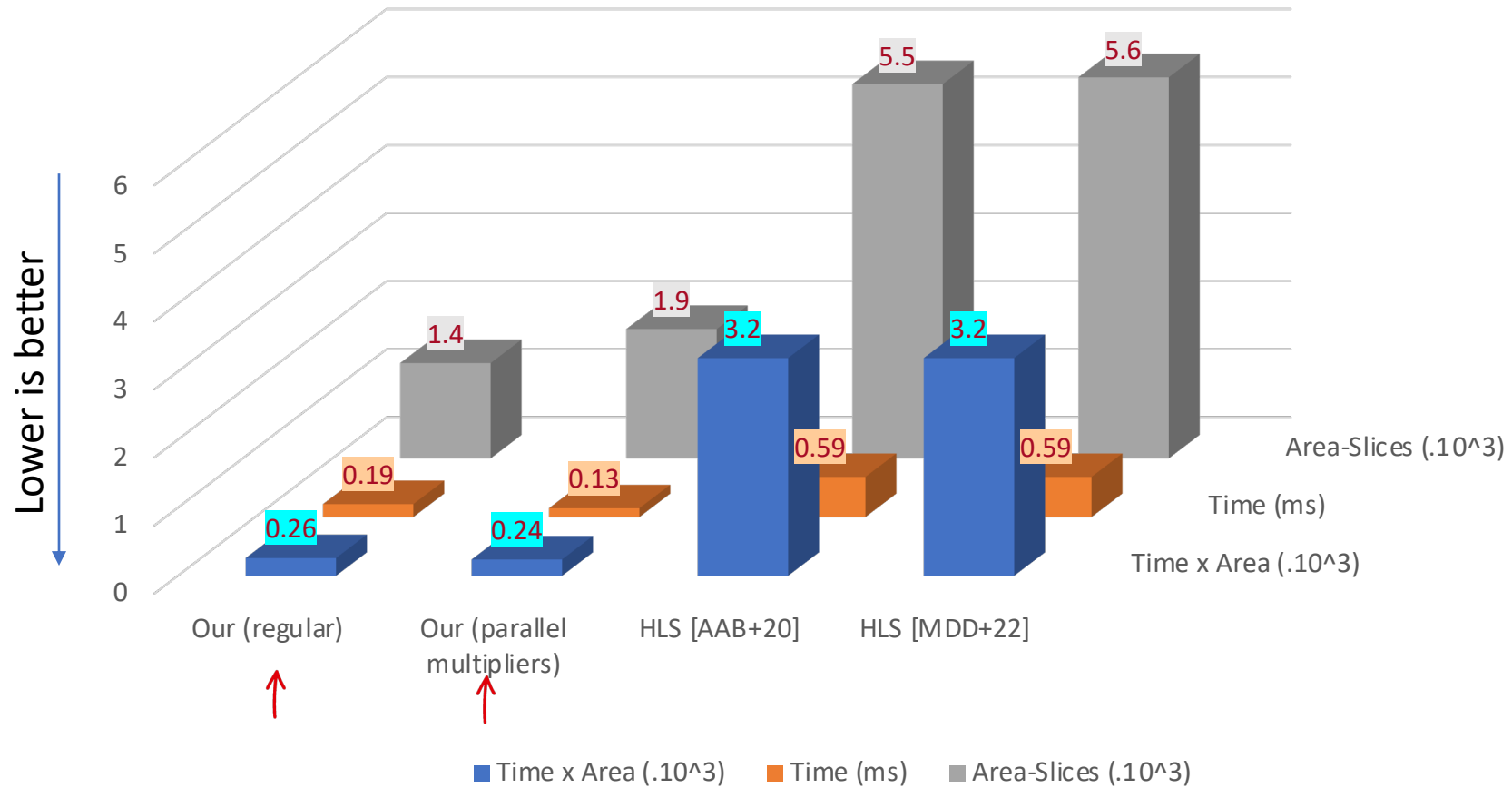
Outputs: Ciphertext (u_{out} , v_{out}), Shared Secret (K_{out}),
Hashed message (d_{out})

1. $(u_{out}, v_{out}) = \text{Encrypt}(m_{in})$
2. $d = \text{Hash}(m_{in})$
3. $SS = \text{Hash}(m_{in}, u_{out}, v_{out})$



Encapsulation – Hardware Design Performance Comparison

HQC128



Hardware Implementation of HQC KEM

1. Key Generation

- Fixed weight error vector generation
 - SHAKE256
- Polynomial Multiplication
- Polynomial Addition

3. Decapsulation

- Decrypt
 - Decode
 - Polynomial Subtraction
 - Polynomial Multiplication
- Encapsulation

2. Encapsulation

- Encrypt
 - Encode
 - Polynomial Multiplication
 - Polynomial Addition
 - Fixed weight error vector generation
 - SHAKE256
- Hash Processing
 - SHAKE256

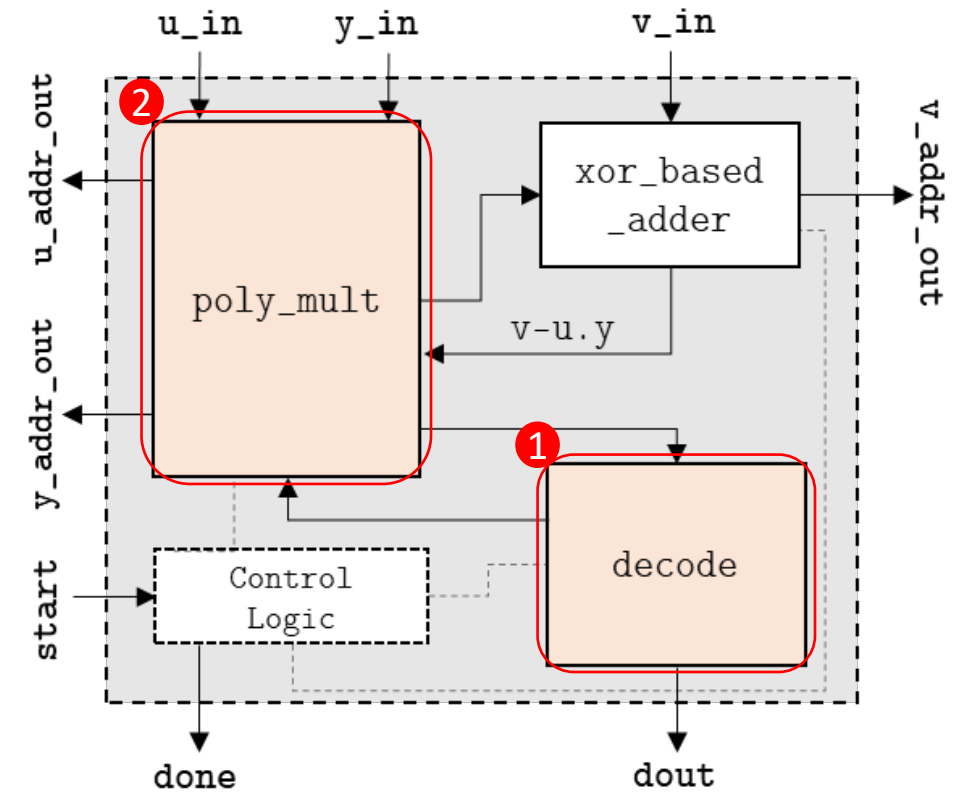
Decrypt - Algorithm and Hardware Design

Inputs: Secret Key (y_{in}), Ciphertext (u_{in}, v_{in})

Outputs: Decoded Message ($dout$)

1. $v_minus_uy = v_{in} - u_{in}.y_{in}$
2. $dout = Decode(v_minus_uy)$

- Decode consists of Reed-Muller and Reed-Solomon decoding



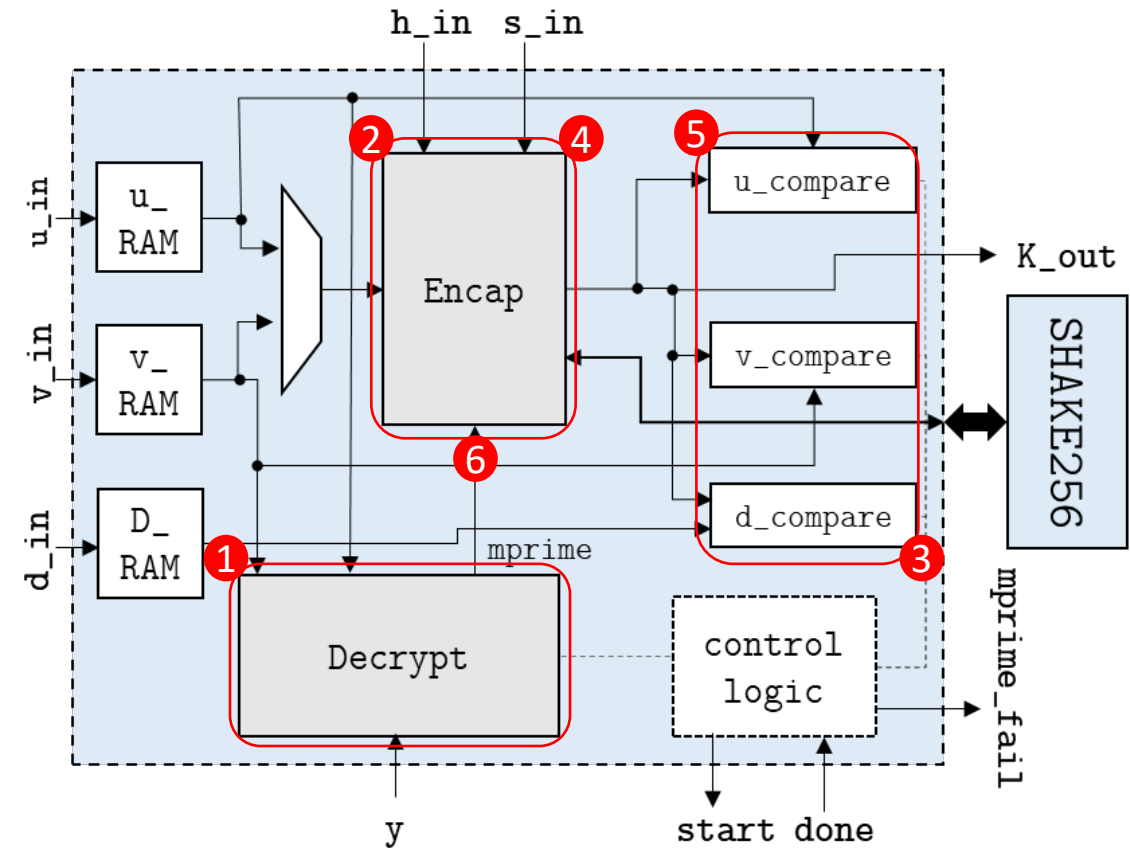
Decapsulation

Inputs: Public Key(h_in , s_in), Secret Key (y_in), Ciphertext (u_in , v_in), Hashed message (d_in)

Outputs: Shared Secret (K_out)

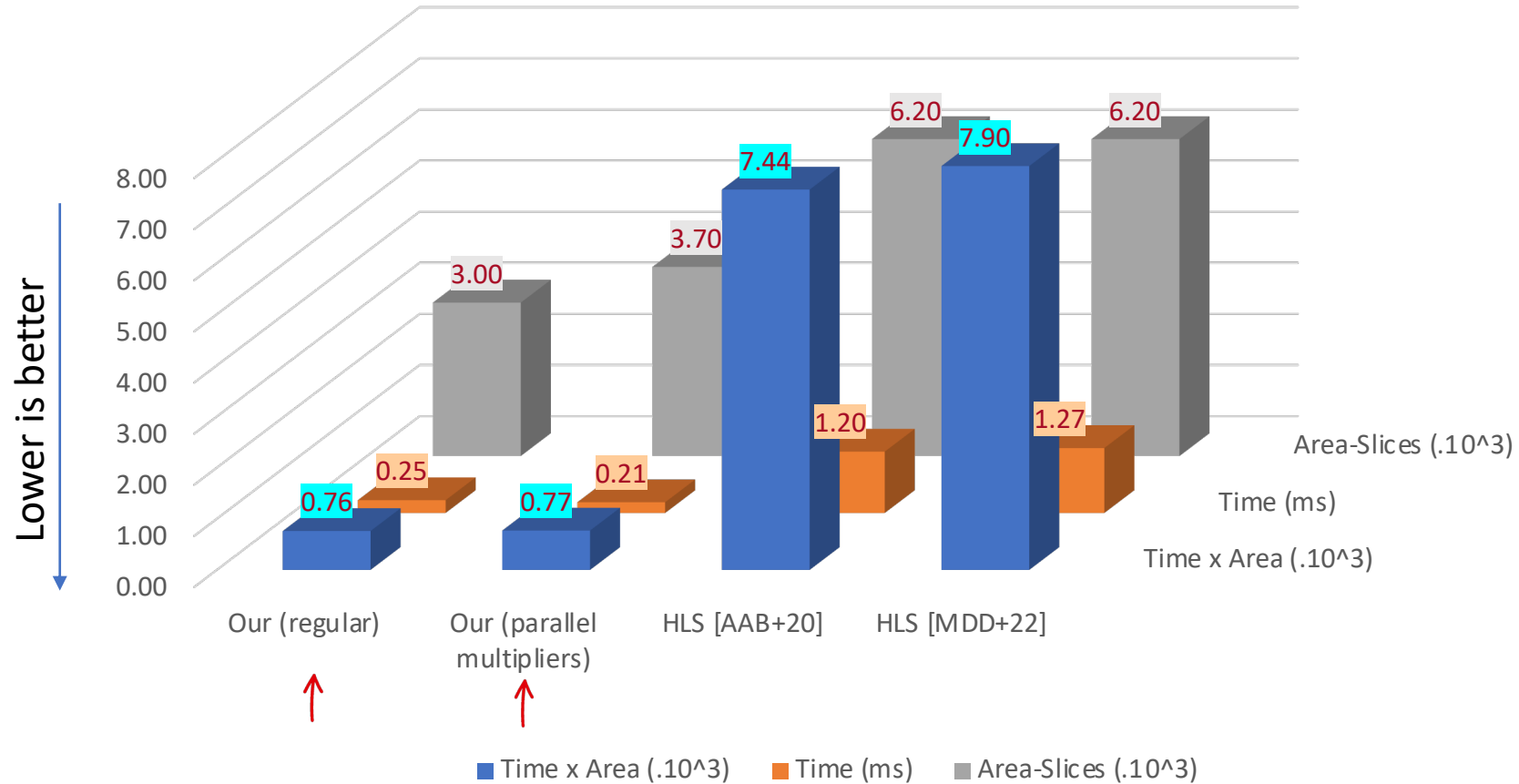
1. $m' = \text{Decode}(v_in - u_in.y_in)$
2. $d' = \text{Hash}(m')$
3. Verify $d_in == d'$?
4. $(u', v') = \text{Encrypt}(m')$
5. Verify $(u', v') == (u_in, v_in)$?
6. $K_out = \text{Hash}(m', u_in, v_in)$

- Decapsulation uses a tweaked version of Encapsulation module where u_prime and v_prime could be swapped with u_in and v_in .



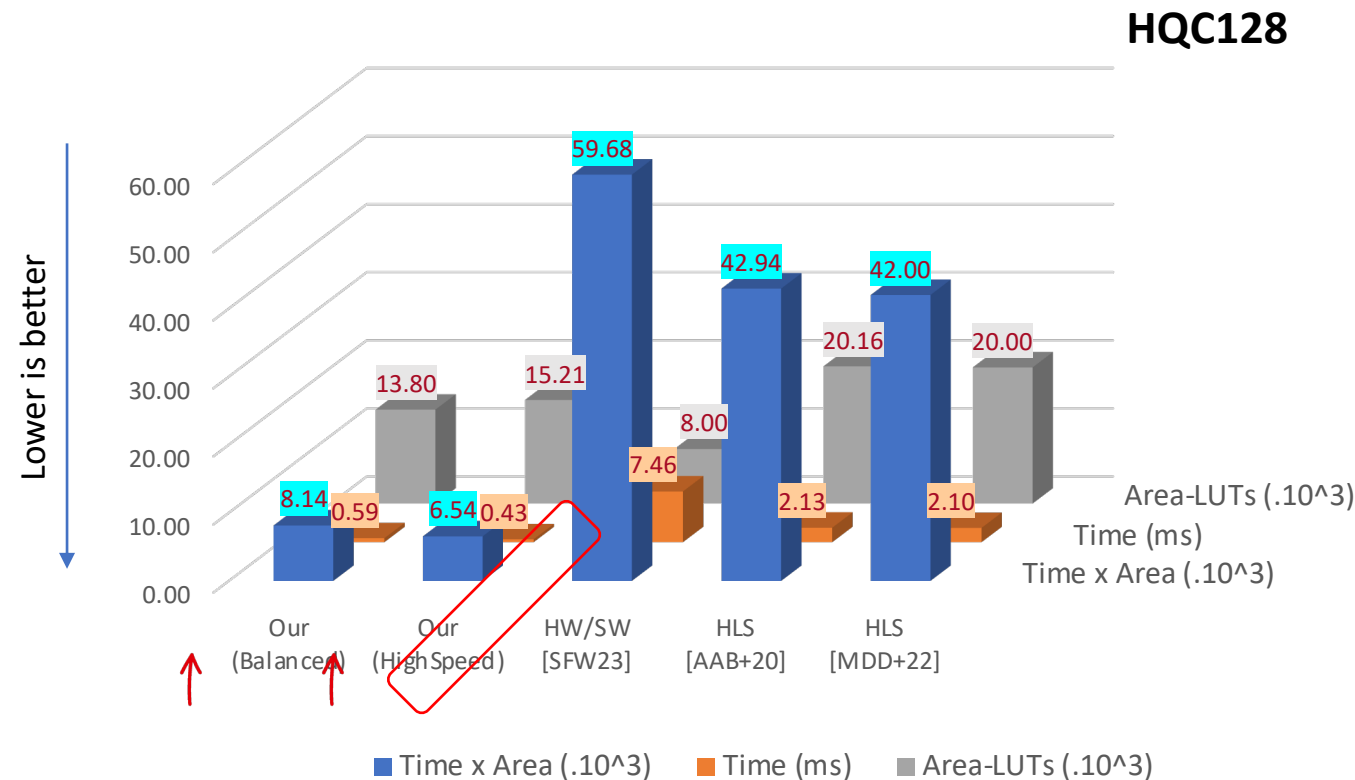
Decapsulation – Hardware Design Performance Comparison

HQC128



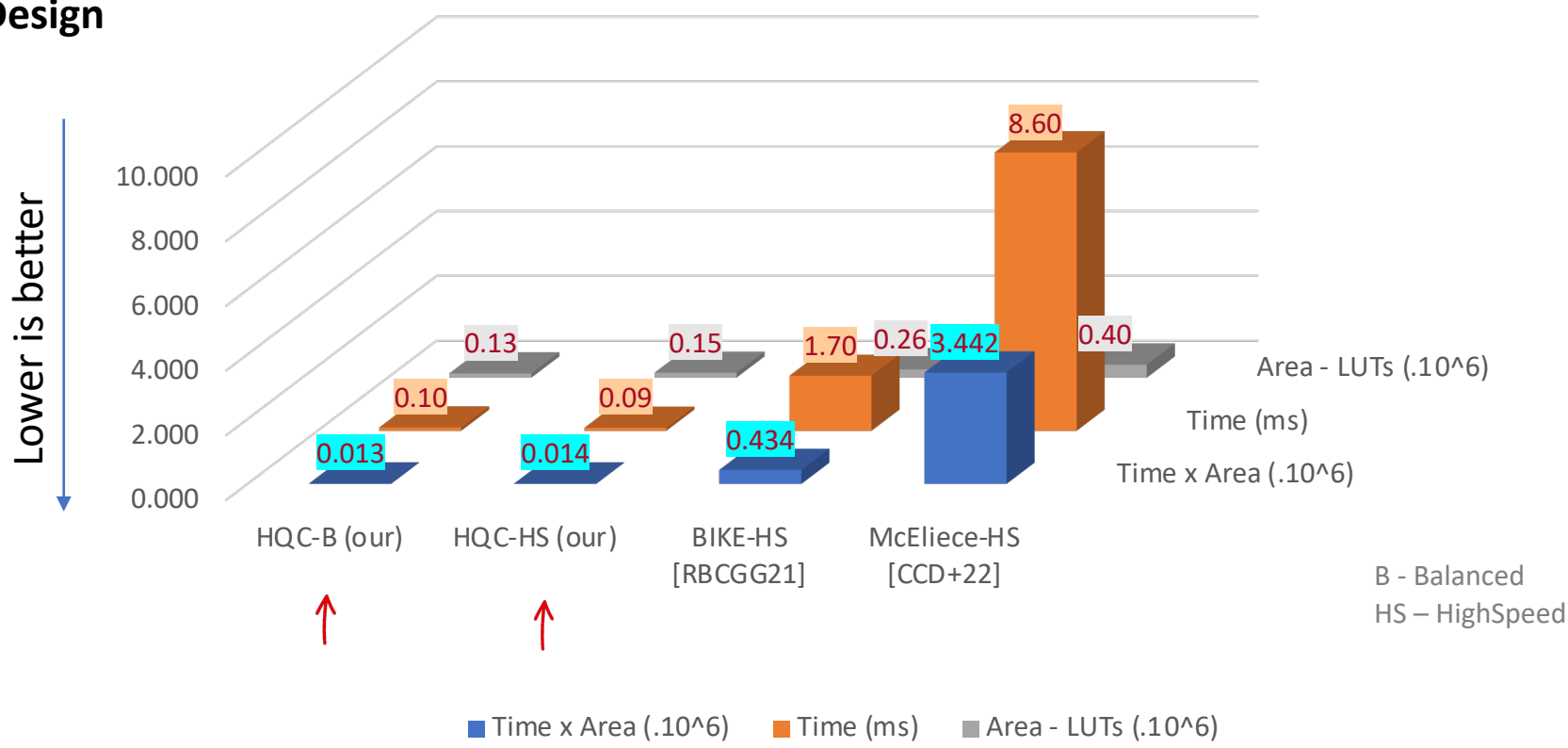
Joint Design – Comparison with other HQC Designs

- Joint Design – combines KeyGen, Encap, and Decap in to one.
- Shared module among different primitives.
 - SHAKE256
 - Polynomial Multiplication
 - Encapsulation
 - Polynomial Addition
- Time = KeyGen + Encap + Decap.



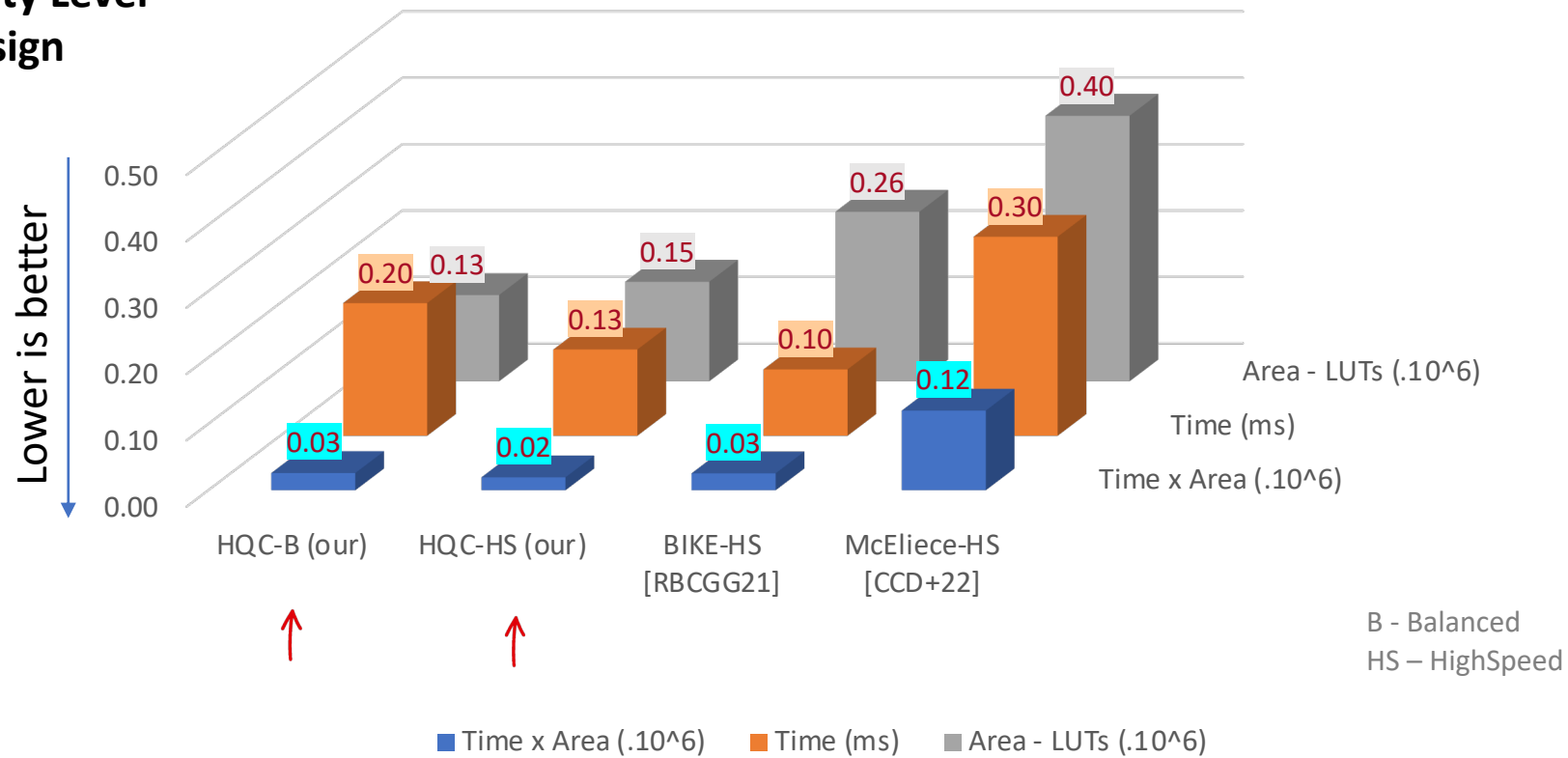
Comparison with other Code Based Schemes - Key Generation

128-bit Security Level
Hardware Design



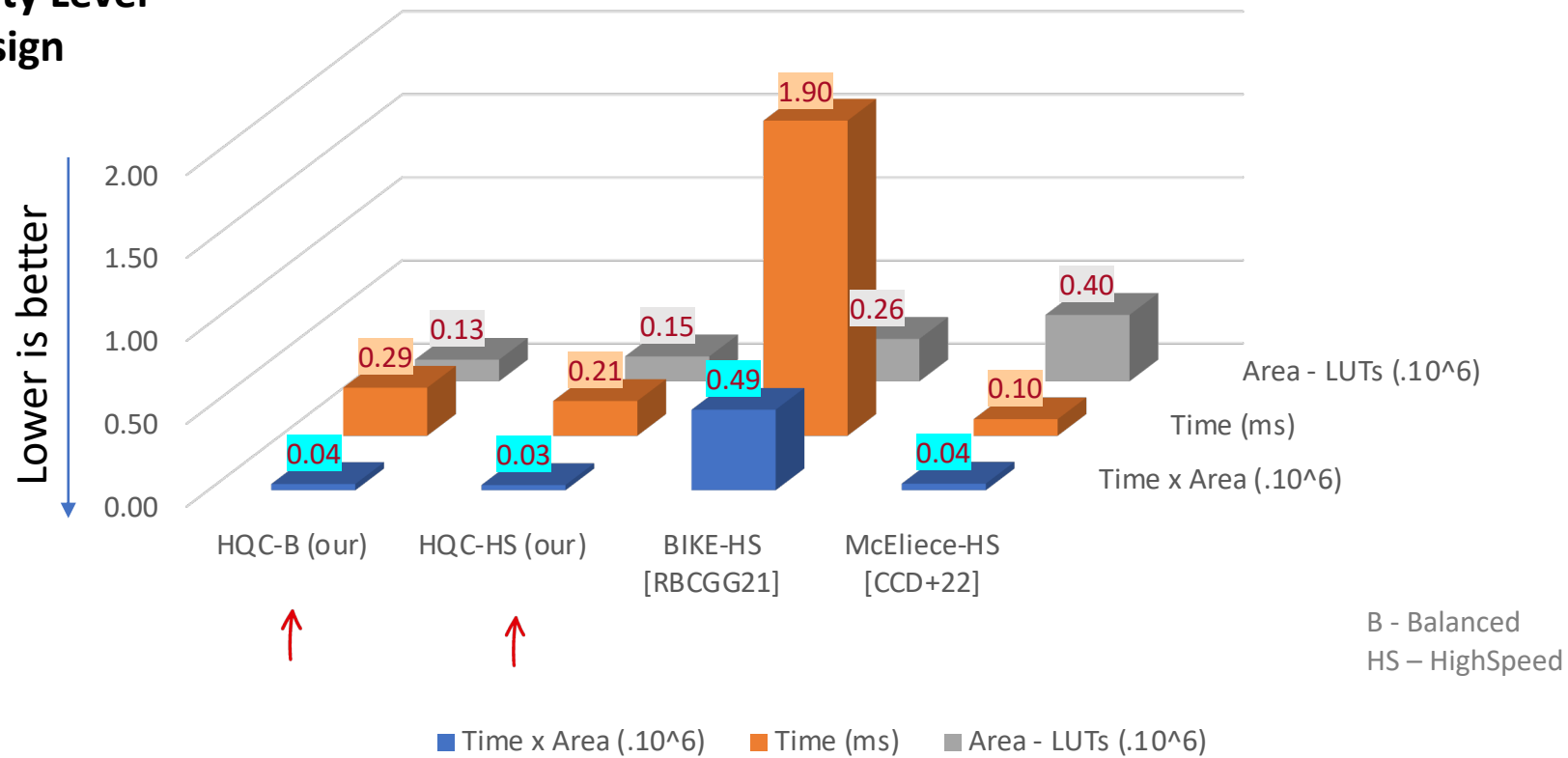
Comparison with other Code Based Schemes - Encapsulation

128-bit Security Level
Hardware Design



Comparison with other Code Based Schemes - Decapsulation

128-bit Security Level
Hardware Design



Summary

- First (hand-tailored) hardware implementations of HQC Key Encapsulation Mechanism parameterizable at compile-time across all parameter sets.
- HQC can be a competitive candidate when optimized hardware is developed.

References

- [AAB+20]** Carlos Aguilar Melchor, Nicolas Aragon, Slim Bettaieb, Loïc Bidoux, Olivier Blazy, Jean-Christophe Deneuville, Philippe Gaborit, Edoardo Persichetti, Gilles Zémor, and Jurjen Bos. HQC. Technical report, National Institute of Standards and Technology, 2020. available at https://pqc-hqc.org/doc/hqc-specification_2021-06-06.pdf
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Summary

- First (hand-tailored) hardware implementations of HQC Key Encapsulation Mechanism parameterizable at compile-time across all parameter sets.
- HQC can be a competitive candidate when optimized hardware is developed.

Thank you!
Questions?



Link to the code base



caslab.io